# Finite fields (Chap. 11)

Some of our new rings IF, [x]/(f(x)) actually turn out to be new (finite) fields IF, with 9= ptg(+).

EXAMPLES Let's write the multiplication tables for

 $R = F_2(x)/(x^2+x+1)$  versus  $R = F_2(x)/(x^2+x)$ 

Livreducible in F. [x] (Why?)

renaming  $\alpha:=\overline{x}$ , so  $R_1 = \text{span}_{F_2} \{1, \alpha\}$   $= \{0, 1, \alpha, \alpha+1\}$ with  $\alpha^2 + \alpha + 1 = 0$ i.e.  $\alpha^2 = \alpha + 1$ 

	×	0	1	X	X+l	
	0	0	0	0	0	
•	1	0	1	×	oct	A
•	K		×		1	field v
	X+I	0	outl	1	d	

(Why?)

# Irreducibility for f(x) is the key:

PROPOSITION: If f(x) is irreducible in F(x) for Fafield, then IF(x)/(f(x)) is also a field. In partoular, if IF=IFp has p elements then IHx]/(f(x)) is a new field with pd elements, where d:= deg(f(x)).

#### EXAMPLES

(1) X71 in IRIX] is irreducible, IR[x]/(x+1) is a field, namely = spange (1, x) our disquised version of the field C= span 1R 11, i] who i2+1=0

(2) x2+x+1 in F2[x] is irreducible, of degree 2, 50  $\mathbb{F}_4 := \mathbb{F}_2[x]/(x^2+x+1) = \{0,1,\alpha,\alpha+1\}$ is a field with  $p^d = 2^2 = 4$  elements

WARNING ?: #4 # 2/4 = {0,1,0,001} = {0,1,2,3}

(3) x4+x+1 in F\_[x] is also irreducible (seen on HW) and has degree 4, so  $F_{16}$ :=  $F_{2}[x](x^{\gamma}+x+1)$  is a field with  $2^{\gamma}=16$  elements = span [1, t, 8, 83] where 8:= x his 14/1=0 i.e. YY = Y+1 = { 0,1,8,841, 841, 848, 84841, 83, 83+1, 83+8, 83+8+1, 83+82+1, 83+828+1)

proof of PROPUSITION:

To see IF(x)/(f(x)) is a field for f(x) irreducible, consider any  $\overline{g(x)} \neq \overline{0}$  in F[x]/(f(x)), and find g(x) as follows. Represent g(x) by some polynomial g(x) with deg(g) < d = deg(f), and then GCD(g(x), f(x)) = 1 since  $g(x) \neq 0$ and fis irreducible. Hence 1 = a(x)f(x) + b(x)q(x) = f(x)and  $\overline{1} = \overline{b(x)} \, \overline{g(x)} \, \text{in } F[x]/(f(x)),$ that is b(x) = g(x) -1.

When 
$$F = F_p$$
, then we know we have a bijection

$$(F_p)^d \longrightarrow F_p(x)/(f(x))$$

$$\begin{bmatrix} c_0 \\ c_1 \\ \vdots \\ c_{d-1} \end{bmatrix} \longrightarrow c_0 + c_1 \overline{x} + c_1 \overline{x}^2 + ... + c_{d-1} \overline{x}^{d-1}$$
so #  $F_p(x)/(f(x)) = \#(F_p)^d = p^d$ 

EXAMPLE (uside 
$$f_{ij} = f_{ij} \times 1/(x^{4} + x + 1)$$
)

with  $Y := \bar{x}$ , who is  $(Y^{2})^{-1}$ ?

 $Y^{2} = \bar{x}^{2}$ , so compute  $Y$  in (extended) finction

 $CD(x^{2}, x^{4} + x + 1) = CD(x + 1, x^{2}) = 1$ 
 $x^{2} = x^{2} = x^{2}$ 

REMARK Although not obvious, any two finite fields  $F_g$  and  $F_g$  having the same size g will be isomorphic, meaning there is a bijection  $F_g$  with  $f(\alpha + \beta) = f(\alpha) + f(\beta)$ .  $f(\alpha \beta) = f(\alpha) \cdot f(\beta)$ .

EXAMPLE The two finite fields of size q=23=8

$$F_8 = F_2[x]/(x^3+x+1) | F_8' = F_2[x]/(x^3+x^2+1)$$
with  $\alpha := \overline{x}$ 

have an isomorphism, e.g., sending  $\alpha \mapsto \beta^3$ .

Check  $\beta^3$  is a root of  $x^3+x+1$  in  $\mathbb{F}'$ :  $(\beta^3)^3 + \beta^3 + 1 = (x^2+1)^3 + (x^2+1)+1$   $= x^4 + x^4 + x^2 + 1 + x^2 + x^4 + x^4$   $= (x^3)^2 + x \cdot x^3 + 1$   $= (x^2+1)^2 + x(x^2+1)+1$   $= x^4 + x^3 + x = x(x^3+x^2+1) = x \cdot 0 = 0$ 

# Primitive voots & primitive polynomials (Chap. 15)

Recall that we showed long ago that if a ring R had units R = [neR: u has a mult. inverse in R] fruite of size N, then every u&Rx had uN=1.

### COROLLARY

In a finite field  $\mathbb{F}_q$  with q elements, every  $\alpha \neq 0$  has  $\alpha^{q-1} = 1$ proof: Fg is a finite ring and # Fg = # (Fg-{0}) = q-1.

DEF'N: Call the smallest power N=1,2,3,... for which  $\alpha=1$  the order of  $\alpha$  in Fig.

Call a a primitive element in II; if it has the largest possible order, namely 9-1.

PEMARK: We'll need primitive voots in Ity later to build Reed-Solomon codes with parameters [9-1, t, q-t] for t = g-1. EXAMPLES

(1) E	= 2/72			has power table				
(1)	Power	1	2	3	4	5	6	
~	1	1	1	1	1	1	1	
	2	2	4	1	2	4	1	
	3	3	2	6	4	5	1	
	4	4	2	1	4	2	1	
,	5	5	4	6	2	3	l	
	6	6	1	6	1	6	1	

(2) In 
$$F_{\gamma} = F_{\Sigma}(x)/(x^{2}x+1)$$
,  $\alpha = \overline{x}$  is primitive since  $\alpha^{2} = \alpha + 1 + 1$ ,  $\alpha = \overline{x}$  is primitive of  $\alpha = 3 - 9 - 1$ .

Also  $\alpha + 1$  is primitive, since  $(\alpha + 1)^{2} = \alpha + 1$ .

(3) In 
$$F_{16} = F_{2}[x]/(x^{q}+x+1)$$
,  $Y = x$  is primitive, that is, of order  $15 = q-1$  but  $Y^{3} = x^{3}$  is not primitive,  $Y^{3} = x^{3} = 1$  since  $(Y^{3})^{5} = Y^{15} = 1$  so  $Y^{3}$  has order  $\leq 5$ , not  $q-1 = 16-1 = 15$ .

Do there exist primitive roots in every finite field Fg? (Yes.) Q: How to knd them?

To answer these, start with some simple properties of order:

- (i) If  $\alpha \in \mathbb{F}^{\times}$  has order d, then  $d=1 \Leftrightarrow dN$
- (ii) Any power at of a has order dividing d (= order)
- (iti) If  $\alpha \in \mathbb{F}^{\times}$  has order d=ef, then  $\alpha^f$  has order e.
- (iv) If  $\alpha_1, \alpha_2 \in \mathbb{F}^{\times}$  have orders  $d_1, d_2$ with GOD(d, d)=1 then ofoly has order d, d2.

sproof:

(i): Certainly if d) N, say N=de, then  $\alpha = \alpha = (\alpha^d)^e = 1^e = 1$ 

On the other hand, if df N then N= de+r with 1≤r≤d-1 so  $\alpha^N = \alpha^{de+r} = (\alpha^d)^e \cdot \alpha^r = 1^e \cdot \alpha^r = \alpha^r \neq 1$  since  $1 \leq r \leq d-1$ .

(ii): If 
$$\alpha^d = 1$$
, then  $(\alpha^k)^d = \alpha^{kd} = (\alpha^a)^k = 1^k = 1$ .

(iii): One has 
$$(\alpha^f)^e = \alpha^{ef} = \alpha^d = 1$$
, and for  $e' < e$  one has  $(\alpha^f)^{e'} = \alpha^{e'} \neq 1$  since  $e' \neq 1 \leq e' \leq d$ .

(iv): 
$$(\alpha_1 \alpha_2)^{N} = 1 \Leftrightarrow \alpha_1^N \cdot \alpha_2^N = 1$$
 $\Rightarrow \alpha_1^N = \alpha_2^N$ 

has order masorder dividing  $d_2$ 

hence it has order dividing  $1 = GD(d_3d_2)$ 
 $\Rightarrow \text{ it is } 1$ .

 $\Rightarrow \alpha_1^N = 1 = \alpha_2^N$ 
 $\Rightarrow \text{ N is a multiple of } d_1 \text{ and } \text{ of } d_2$ 
 $\Rightarrow d_1 d_2 \mid N \mid E$ 

THEOREM Every finite field  $\mathbb{F}_{g}$  has a primitive voot. Proof: Let l := LCM forders of all  $\alpha \in \mathbb{F}_{g}^{\times}$ ?

(e.g. for  $\mathbb{F}_{4} = \{ \}, 1, 2, 3, 9, 5, 6, 2 \}$ orders
1 3 6 3 6 2

50 l = LCM(1,3,6,3,6,2) = 6

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We dain that l= g-1 = # Fg): First note 1 9-1 since every of Fx has 07=1 and so its order divides q-1, and thus so does their LCM Second note 1 29-1 because every  $\alpha \in \mathbb{F}_q^{\times}$ has  $x^{2}=1$  making it a root of  $f(x)=x^{2}-1$ , which cannot have more tran I distinct wots. This proves the claim that l= 9-1. Now we show I some  $\alpha \in \mathbb{F}_q^{\times}$  of order L (=q-1), which would then be a primitive element.

Tactor  $l = p_1^p p_2^p - p_1^p$  into prime powers  $p_i^p$  (for distinct primes). Since l=LCM (orders of XEFx), for each i=1,2=5 there must be some  $xi \in \mathbb{F}_q^{\times}$  with order divisible by  $p_i$ . Then some power of has order exactly Pi.
And then  $a := \alpha_1^{d_1} \alpha_2^{d_2} - \alpha_r^{d_r}$  will have order exactly pipe-pr=l using GOD(pi, pi)=1 repeatedly. 121

So primitive voots exist in Eq,
but how to actually find one?

His slightly tricky—

many elements of Fg are primitive

(in fact, exactly 4(q-1) of them; see §15.8).

So a strategy is to look for them via random search, once we have a quick test for primitivity:

PROPOSITION:  $\alpha \in \mathbb{F}_{x}^{\times}$  is primitive  $(\S16.3)$   $\iff \alpha \neq 1$   $\forall$  primes  $p \mid q-1$  (i.e. if  $q-1=p_{1}^{*} \cdots p_{r}^{*}$  then check  $\alpha \neq 1$   $\forall$  primes  $p \mid q-1$ )

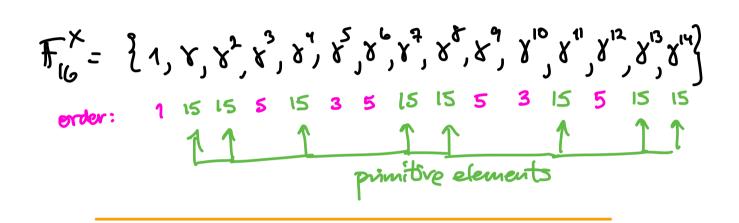
proof:  $\alpha \in \mathbb{F}_{x}^{\times}$  has  $\alpha^{q-1} = 1$ , so its order  $\alpha \mid q-1$ , and  $\alpha \in \mathbb{F}_{x}^{\times}$  has  $\alpha^{q-1} = 1$ , so its order  $\alpha \mid q-1$ ,  $\alpha \in \mathbb{F}_{x}^{\times}$  has  $\alpha \mid q-1$   $\Rightarrow \alpha \mid q$ 

## EXAMPLES

- (1) To check which elements in  $\mathbb{F}_g^{\times}$  are primitive, factor  $q^{-1} = 8^{-1} = 7$  into primes (only one!) and then  $X \in \mathbb{F}_g^{\times}$  is primitive  $\Leftrightarrow \chi = \chi = \chi = \chi = 1$  i.e. the other 6 elements in  $\mathbb{F}_g^{\times}$ -11) are all primitive.
- (2) In  $\mathbb{H}_{16}^{\times}$ , factor  $q = 16 1 = 15 = 3.5^{1}$ and then  $\alpha \in \mathbb{H}_{16}^{\times}$  is primitive  $\Leftrightarrow 1 \neq \alpha = 16 - 1 = 15 = 3.5^{1}$ and then  $\alpha \in \mathbb{H}_{16}^{\times}$  is primitive  $\Leftrightarrow 1 \neq \alpha = 16 - 1 = 15 = 3.5^{1}$ and  $1 \neq \alpha = 3 = 3$

So when we built  $H_{16} = H_2[x]/(x^4+x+1)$ , this is how we would check Y:=x was printitle:  $Y^3 \neq 1$  (since  $H_{16}$  has  $H_2$ -basis  $\{1, Y, \chi^2, \chi^3\}$ )  $Y^5 = Y \cdot \chi^4 = \chi(\chi_{+1}) = \chi^2 + \chi + 1$ 

REMARK: Once we know  $\gamma$  is primitive, the other primitive voots are easier to spot, because they're of the form  $\gamma^i$  for  $i \in 1, 2, ..., q-1$  with GCD(i, q-i)=1:



REMARK: Primitive voots also help us find CRC generator polyomials g(x) in Fp/x] that outch 2-digit errors that are tar apart.

Recall g(x) as a CRC catches such errors up to N digits apart where N is smallest such that  $g(x) \mid x^N - 1$  in Fp[x].

PROPOSITION: Let  $(x) \in \mathbb{F}_0[x]$  be irreducible, so that  $\mathbb{F}_q = \mathbb{F}_0[x]/(g(x))$  is a field of size  $q = p^d$  where d = deg(g(x)).

Then the smallest N for which  $g(x)|x^{-1}$  is the order of  $\alpha := x$  in this field f(x).

In particular, N divides q-1=pd-1, with equality N=pd-1 ⇒ x is a primitive root.

In this case, one calls g(x) a primitive (wedneible) polynomial in Fp(x).

proof: 
$$g(x) \mid X^N - 1$$
 in  $\mathbb{F}_p[x]$ 
 $\Rightarrow \widehat{\chi}^N - \overline{1} = \overline{0}$  in  $\mathbb{F}_p[x]/(g(x))(=:\mathbb{F}_g)$ 
 $\Rightarrow \overline{1} = \overline{\chi}^N = \alpha^N$ 

Hence the smallest such  $N$  is the order of  $\alpha$ 

upstot: Primitive polynomials g(x) in Fp(x) make very good choices for CRC's catching 2-digit errors.

#### EXAMPLES

- (1) In  $F_{16} = \frac{F_2[x]}{(x^4x+1)}$ , we checked g(x)  $x = \overline{x}$  was a primitive root, so  $g(x) = x^4 + x + 1$  in  $F_2[x]$  is a primitive polynomial, and used as a CRC will catch 2-bit errors up to N = 16 1 = 15 bits apart.
- (2) Garrett mentions  $g(x)=x^{15}+x+1 \in \mathbb{F}_2[x]$  as being a primitive polynomial, so as a CRC it will catch 2-bit errors up to  $N=2^{15}-1=32767$  bits apart!