## Study Guide for Exam 1: Solutions

1. (a) First, recall ${ }^{1}$ that the disjuction $A \vee B$ of two statements $A$ and $B$ is true if and only if at least one of the statements is true, and the conjunction $A \wedge B$ is true only when both $A$ and $B$ are true.
Now, let $R$ be a true statement, while $P$ and $Q$ are false. Then the disjunction $\underbrace{(P \wedge Q)}_{F} \vee \underbrace{R}_{T}$ is true, but the conjunction $\underbrace{P}_{F} \wedge \underbrace{(Q \vee R)}_{T}$ is false.
Remark. Compare the given formulas with the ones given in exercise 1.14(c)-(f).
(b) A quick reminder: the implication ${ }^{2} A \Rightarrow B$ is false only in the case when $A$ is a true statement and $B$ is false.

## Solution 1.

Comparing the truth tables for both of the given statements, we conclude that they are equivalent:

|  |  | $P$ | $Q$ | $R$ | $P \wedge Q$ | $(P \wedge Q) \Rightarrow R$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | T | T | T | T | T |  |
|  |  | T | T | F | T | F |  |
|  |  | T | F | T | F | T |  |
|  |  | T | F | F | F | T |  |
|  |  | F | T | T | F | T |  |
|  |  | F | T | F | F | T |  |
|  |  | F | F | T | F | T |  |
|  |  | F | F | F | F | T |  |
| $P$ | $Q$ | $R$ | $P \Rightarrow R$ |  | $Q \Rightarrow R$ | (P ${ }^{\text {a }}$ ( $) \vee(Q \Rightarrow R)$ |  |
| T | T | T | T | T | T | T |  |
| T | T | F | F | F | F | F |  |
| T | F | T |  | T | T | T |  |
| T | F | F | F | F | T | T |  |
| F | T | T |  | T | T | T |  |
| F | T | F |  | T | F | T |  |
| F | F | T |  | T | T | T |  |
| F | F | F |  | , | T | T |  |

## Solution 2.

Since for any two statements $A$ and $B$, the implication $A \Rightarrow B$ is equivalent to the disjunction $\sim A \vee B$, then

$$
(P \wedge Q) \Rightarrow R=\sim(P \wedge Q) \vee R \stackrel{\text { De Morgan }}{=}(\sim P \vee \sim Q) \vee R .
$$

Similarly,

$$
(P \Rightarrow R) \vee(Q \Rightarrow R)=(\sim P \vee R) \vee(\sim Q \vee R)=(\sim P \vee \sim Q) \vee R
$$

Thus, the given statements are equivalent.

[^0]2. (a) $\forall a \in \mathbb{R} \forall b \in \mathbb{R}, P(a, b) \Rightarrow\left(P\left(a, \frac{a+b}{2}\right) \wedge P\left(\frac{a+b}{2}, b\right)\right)$.
(b) $\forall x \in \mathbb{R},(P(0, x) \wedge P(x, 1)) \Rightarrow P\left(x^{2}, x\right)$.
(c) $\exists x \in \mathbb{R} \ni P(0, x) \wedge P(x, 1) \wedge \sim P\left(x^{2}, x\right)$.
3. (a) Let $A$ and $B$ be subsets of $\mathbb{R}$. Then the given statement $p$ is, basically, the implication $r \Rightarrow s$, where $r$ is the statement
the intersection $A \cap B$ is infinite,
and the statement $s$ reads
$$
\text { both } A \text { and } B \text { are infinite. }
$$

Recall ${ }^{3}$ that the contrapositive of an implication $r \Rightarrow s$ is the implication $\sim s \Rightarrow \sim r$. In our case
$\sim r$ : the intersection $A \cap B$ is finite,
$\sim s$ : at least one of the sets $A$ and $B$ is finite.

Hence, the contrapositive of $p$ can be stated as

$$
\text { if at least one of the sets } A \text { and } B \text { is finite, then } A \cap B \text { is finite. }
$$

(b) Let $r$ and $s$ be as above. By definition ${ }^{4}$, the converse of an implication $r \Rightarrow s$ is the implication $s \Rightarrow r$. In our case it reads
if both $A$ and $B$ are infinite sets, then the intersection $A \cap B$ is infinite.
(c) The converse of $p$ is false. Here is a counterexample: let $A$ be the set of all even integers and $B$ be the set of all odd integers. Then both $A$ and $B$ are infinite, but $A \cap B=\varnothing$ is a finite set.
Remark. Make sure to know the difference between contrapositive, converse, inverse and negation of an implication $r \Rightarrow s$ :

|  | Formula | Remarks |
| :--- | :---: | :--- |
| Contrapositive | $\sim s \Rightarrow \sim r$ | $\begin{array}{l}\text { It is logically equivalent to the orig- } \\ \text { inal implication } r \Rightarrow s .\end{array}$ |
| Converse | $s \Rightarrow r$ | $\begin{array}{l}\text { Be careful. It is not equivalent to } \\ \text { the original statement } r \Rightarrow s .\end{array}$ |
| Inverse | $\sim r \Rightarrow \sim s$ | $\begin{array}{l}\text { It is logically equivalent to the con- } \\ \text { verse. }\end{array}$ |
| Negation | $r \wedge \sim s$ | $\begin{array}{l}\text { The negation of an implication is } \\ \text { not an implication. } \\ \text { Exercise. Let } r \text { be the statement } \\ \text { " }\end{array}$ |
| you give a mouse a cookie", and |  |  |
| denote the statement "he will ask |  |  |
| for a glass of milk". Then the im- |  |  |
| plication $r \Rightarrow s$ is a popular saying. |  |  |
| How does the negation of this im- |  |  |
| plication look like? |  |  |$]$

[^1]
[^0]:    ${ }^{1}$ Chapter 1, pp.3-4
    ${ }^{2}$ Chapter 1, p. 5

[^1]:    ${ }^{3}$ Chapter 1, p. 21
    ${ }^{4}$ Chapter 1, p. 22

